

# Quantum Circuits

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<http://aggregate.org/hankd/>

# Why Quantum Algorithms?

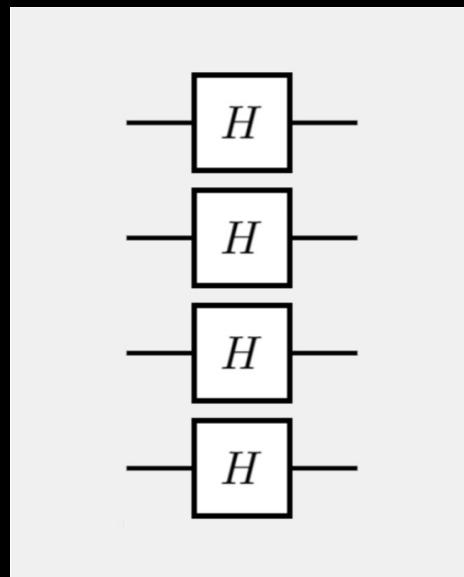
- Find a solution faster:
  - Exponentially parallel execution
  - Quantum operations reduce  $O()$  complexity
- Reduce memory size required:  
Holding  $2^n$   $n$ -bit values in  $n$  qubits
- Reduce power consumed per computation:
  - Parallel computation without parallel HW
  - Reduced  $O()$  complexity reduces operations

# Why Quantum Algorithms?

- Find a solution faster:
  - Exponentially parallel execution
  - Quantum operations reduce  $O()$  complexity
- Reduce memory size required:  
 $Holding 2^n n\text{-bit values in } n \text{ qubits}$
- Reduce power consumed per computation:
  - Parallel computation without parallel HW
  - Reduced  $O()$  complexity, fewer operations

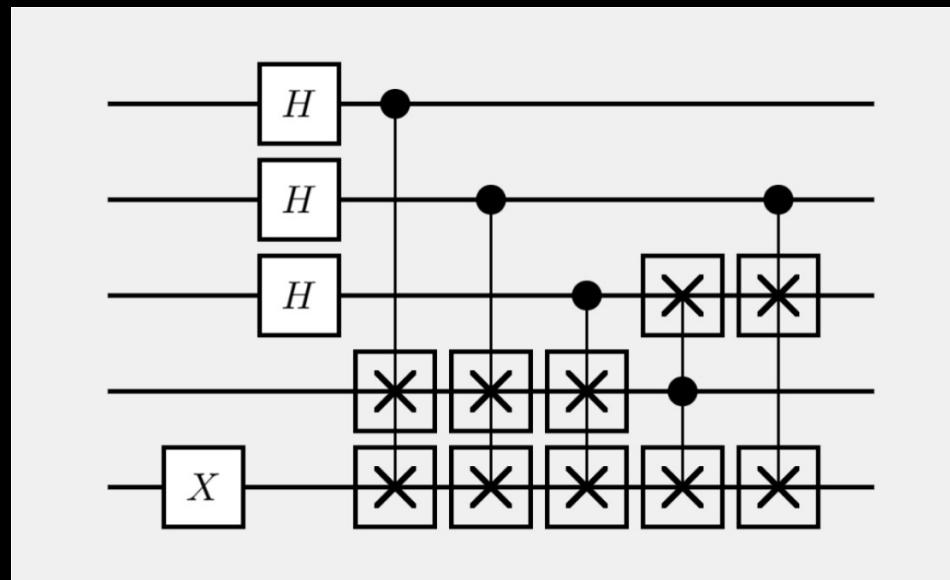
# Parallel Evaluation

- Perform operation on all data values
  - Measure a randomly-selected result
  - **MuqcsCraft Random 4-bit value**



# Parallel Evaluation

- Perform operation on all data values
  - Report a randomly-selected result
  - **MuqcsCraft FA**



# Query Model of Computation

- Input isn't data, but a function
  - Goal is finding a property of the function by making queries against it
  - Function is treated as an **oracle**
  - Function is computable using binary logic
- The function is mapped into a unitary gate  $U_f$ 
  - $U_f$  inputs  $|x\rangle|y\rangle$  and outputs  $|x\rangle|y^f(x)\rangle$
  - $f(x)$  can be 0 or 1; *not necessarily* entangled
  - $U_f$  does not have to be simple...

# Deutsch's Problem

- Data isn't the input; a function  $f(a)$  is
- There are 4 possible functions of one qubit:

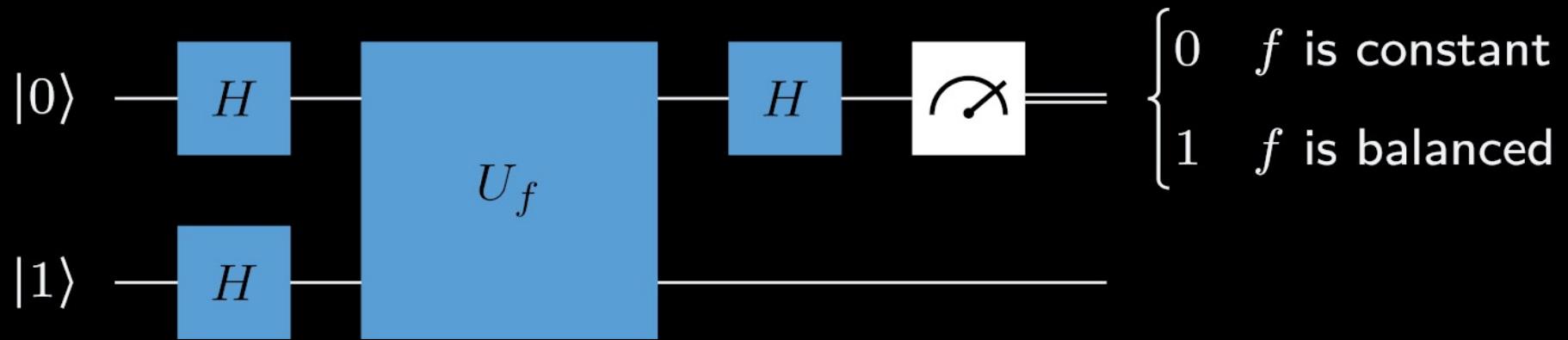
$a$	$f_0(a)$	$f_1(a)$	$f_2(a)$	$f_3(a)$
0	0	0	1	1
1	0	1	0	1

- A function  $f(a)$  can be:
  - **Constant**: always same output:  $f_0(a)$ ,  $f_3(a)$
  - **Balanced**:  $\{0,1\}$  equiprobable:  $f_1(a)$ ,  $f_2(a)$
  - Neither: *well, not in this case...*

# Conventional Solution

- Query  $f(a)$  twice:  $f(0)$  and  $f(1)$   
*that is testing all possible values...*
- Is  $f(0) == f(1)$ ?
  - Yes: function is **Constant**
  - No: function is **Balanced**

# Deutsch's Algorithm



- $X$  input passes through  $U_f$  unchanged (generally necessary to make  $U_f$  reversible)
- $Y$  input is **XORed** with  $f(X)$  in  $U_f$ , but  $Y$  is  $180^\circ$  out of phase with  $X$  (due to  $|1\rangle$ )
- Final  $H(X)$  completes **phase kickback**

# Deutsch's Algorithm

$a$	$f_0(a)$	$f_1(a)$	$f_2(a)$	$f_3(a)$
0	0	0	1	1
1	0	1	0	1

- $U_{f0}$  inputs  $|x\rangle|y\rangle$  and outputs  $|x\rangle|y^{\wedge}f_0(x)\rangle$ 
  - $f_0(x)$  is 0
  - $U_{f0}$  is  $|x\rangle|y^{\wedge}0\rangle$  which is just  $|x\rangle|y\rangle$
  - Thus,  $U_{f0}$  is no gates at all!
- **MuqcsCraft Uf0**, entangled **MuqcsCraft Uf0**

# Deutsch's Algorithm

$a$	$f_0(a)$	$f_1(a)$	$f_2(a)$	$f_3(a)$
0	0	0	1	1
1	0	1	0	1

- $U_{f1}$  inputs  $|x\rangle|y\rangle$  and outputs  $|x\rangle|y^{\wedge}f_1(x)\rangle$ 
  - $f_1(x)$  is  $x$
  - $U_{f1}$  is  $|x\rangle|y^{\wedge}x\rangle$
  - Thus,  $U_{f1}$  is a **CNOT**  $x, y$  gate
- **MuqcsCraft Uf1**

# Deutsch's Algorithm

$a$	$f_0(a)$	$f_1(a)$	$f_2(a)$	$f_3(a)$
0	0	0	1	1
1	0	1	0	1

- $U_{f_2}$  inputs  $|x\rangle|y\rangle$  and outputs  $|x\rangle|y^{\wedge}f_2(x)\rangle$ 
  - $f_2(x)$  is  $\sim x$
  - $U_{f_2}$  is  $|x\rangle|y^{\wedge}\sim x\rangle$  which is just  $|x\rangle|\sim(y^{\wedge}x)\rangle$
  - Thus,  $U_{f_2}$  is a **CNOT x, y** , **NOT y** gate
- **MuqcsCraft Uf2**

# Deutsch's Algorithm

$a$	$f_0(a)$	$f_1(a)$	$f_2(a)$	$f_3(a)$
0	0	0	1	1
1	0	1	0	1

- $U_{f_3}$  inputs  $|x\rangle|y\rangle$  and outputs  $|x\rangle|y^{\wedge}f_3(x)\rangle$ 
  - $f_3(x)$  is 1
  - $U_{f_3}$  is  $|x\rangle|y^{\wedge}1\rangle$  which is just  $|x\rangle|\sim y\rangle$
  - Thus,  $U_{f_3}$  is a **NOT y** gate
- **MuqcsCraft Uf3**

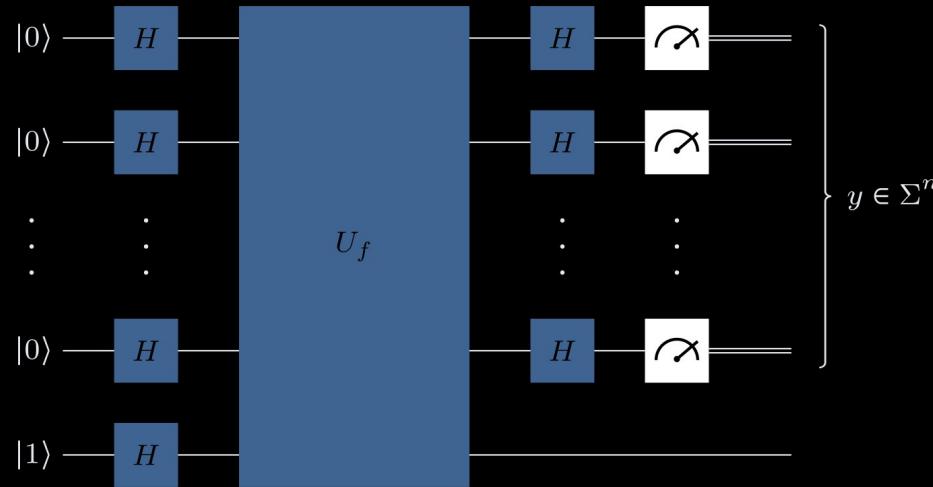
# Deutsch-Jozsa Problem

- Extends Deutsch Problem to operate on a function with  $k$  inputs and one output
- Distinguishes **constant** from **balanced**, but “don’t care” about result if it is neither (i.e., if **promise** is not kept)
- Random functions are unlikely to be either constant or balanced; e.g., **AND**  $X_0, X_1$

# Conventional Solution

- Evaluate  $f()$  for 2 to  $2^{k-1}+1$  random inputs
- If any two evals don't return the same value, it must be balanced and can stop early
- To get a statistical answer, can stop after  $n$  evaluations that were all the same
  - If  $f()$  is constant, answer is correct
  - If  $f()$  is balanced, probability of error is  $2^{-n+1}$

# Deutsch-Jozsa Algorithm



- Output 0 for constant, 1 for balanced, but 1 is if **ANY** measurement is 1 (**OR** reduction)
- **OR** reduction takes  $O(k)$  operations, but  $U_f$  is evaluated just once, with  $2^k$  parallelism

# Deutsch-Jozsa Algorithm

X1	X0	0	NOT X0	AND
0	0	0	1	0
0	1	0	0	0
1	0	0	1	0
1	1	0	0	1

- $f(X1, X0) = 0$  is constant: **MuqcsCraft DJ0**
- $f(X1, X0) = \sim X1$  is balanced: **MuqcsCraft DJNOT**
- $AND(X1, X0)$  is neither: **MuqcsCraft DJAND**

# Bernstein-Vazirani Problem

- Sometimes called **Fourier sampling** problem
- Given the promise that there exists a vector  $s$  such that  $f(x) = s \cdot x$  for all  $x$ , find  $s$

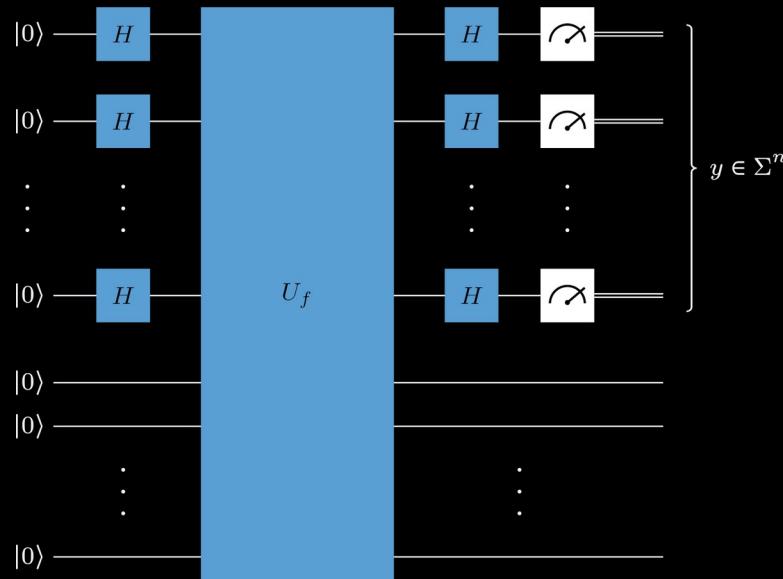
Note that:  $s \cdot x = (s_0 \& x_0) \wedge (s_1 \& x_1) \dots$

- Uses Deutsch-Jozsa Algorithm
  - Measurements are weights for  $s$
  - No conventional postprocessing (no ANY)
  - For example: **MuqcsCraft DJNOT**

# Simon's Problem

- For a function with  $n$  inputs and  $m$  outputs
- Find  $v$  such that  $f(x) == f(y)$  implies either:
  - $x^s == y$
  - $x == y$  (in which case  $s = 0^n$ )
- Requires promise that  $s$  exists...

# Simon's Algorithm



- The  $n$  inputs are at the top and  $m$  outputs are at the bottom – without phase kickback
- Measurement does not directly give s...

# Simon's Algorithm

- Each run of the quantum algorithm gives a randomly-selected  $n$ -bit vector  $y$ ; collect these into a binary matrix  $M$  with  $n$  columns and  $k$  rows (one for each quantum algorithm run)
- $M s$ , for a column vector of  $s$ , should equal 0; thus, classical Gaussian elimination can be used to solve for  $s$
- Classical queries only eliminate one possible  $s$  per pair of queries producing different values, so  $\geq 2^{n/2-1} - 1$  **queries** are generally needed

# Qiskit

- **Qiskit** is the most popular quantum circuit SDK
  - Open source software from IBM Quantum
  - Python, but includes hardware interfaces
  - Qiskit ecosystem collects related projects

<https://www.ibm.com/quantum/ecosystem>

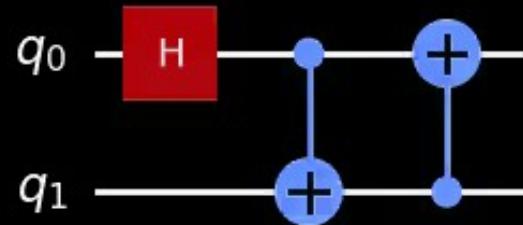
- Install Qiskit on a local machine

<https://quantum.cloud.ibm.com/docs/en/guides/install-qiskit#local>

- IBM Quantum Experience was drag-and-drop, WWW versions are now Jupyter Notebooks

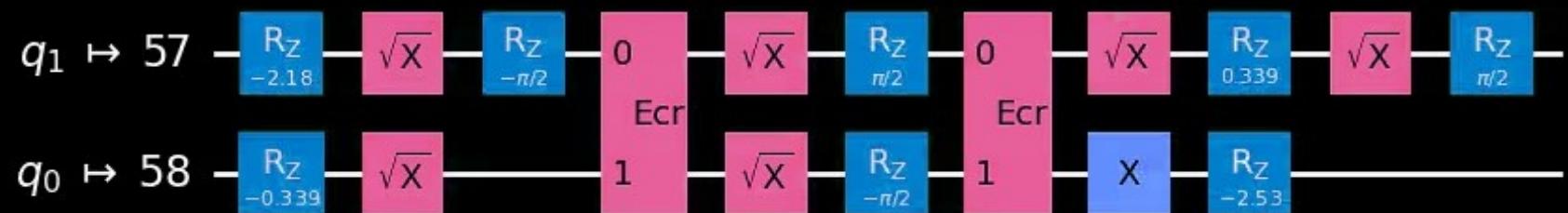
<https://quantum.cloud.ibm.com/docs/en/guides/online-lab-environments>

# Transpilation



- **Transpilation** is compilation
  - Allocating specific qubits
  - Translating to supported primitive operations
  - Applies optimization and scheduling passes

Global Phase:  $2\pi$



# Qiskit philosophy

- Qiskit is a library in conventional python code
- Qiskit constructs a data structure representing each quantum circuit and calls functions to perform actions on that data structure
- IBM's Qiskit in the classroom  
<https://quantum.cloud.ibm.com/learning/en/modules/quantum-mechanics/get-started-with-qiskit>
- Let's look at the “Hello World” example  
<https://quantum.cloud.ibm.com/docs/en/guides/hello-world>