

# Parallel Bit Pattern (part 1)

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**Hank Dietz**

<http://aggregate.org/hankd/>

# To pbits And Beyond

- In part 0, we explained how an AoB can represent an entangled superposition... and how layers up to RE avoid bit-level work
- The layers here are the programming interface presented by a **C++ library**
  - `pbit` – a dynamically managed RE reference
  - `pint` – variable precision, signed/unsigned
  - `pfloat` – variable precision (accuracy)

# The C++ Library: `pbp.h`

- Any C++ program can `#include <pbp.h>`
  - **Warning:** research-grade implementation!
  - Defines `pbit`, `pint`, and `pfloat` classes
  - These classes function like built-in types with operators overloaded to handle them
- Library code lowers to optimized AoB ops
  - Currently, entirely done at runtime (like qiskit)
  - Actual execution is implicitly triggered when needed, some versions use lazy evaluation

# The C++ Library Parameters

- Library is highly portable
  - 32/64 bit processors, SWAR support
  - For ESP32, includes C++ subset interpreter
- Key library parameters:
  - **REWAYS**: max ways entangled ( $\leq 32$ )
  - **AOBWAYS**: ways per chunk (5..9)
  - **REREGS**: how many RE registers
  - **AOBREGS**: how many AOB registers
  - **ACSIZE**: size of applicative cache

## Using the PBP library

The PBP library is written as highly portable, self-contained, C++ code. All that is needed to use it is inclusion of the header file with `REWAYS` set to the desired maximum entanglement (default 10).

```
#include "pbp.h"
```

## Sample pint Layer Algorithms

It is easy to compute the square root of an 8-bit number by exhaustive search. For example, `sqrt(169)` will find 13.

```
void pintsqrt(int val) {
    pint a(val); // 8-bit number
    pint b = H(4); // all possible roots
    pint c = (b * b); // square them
    pint d = (c == a); // select answer
    int pos = d.First();
    printf("Square root of %d is %d\n",
        val, pos);
}
```

A less obvious algorithm factors an 8-bit number. Here, possible 4-bit factors are assigned different entanglement channel sets so the multiply produces an 8-way entangled answer rather than 4-way. For example, `factor(143)` will find 11 and 13.

```
#include "pbp.h"
```

```
void pintfactor(int val) {
    pint a(val); // 8-bit number
    pint b = H(4,0x0f); // 4-bit
    pint c = H(4,0xf0); // 4-bit
    pint d = b * c; // multiply 'em
    pint e = (d == a); // which were val?
    pint f = e * b; // zero non-answers
    int spot = f.First(); // factors
    int one = c.Meas(spot);
    int two = b.Meas(spot);
    printf("%d, %d are factors of %d\n",
        one, two, val);
}
```

As above, algorithms written for PBP tend to use abilities that quantum computers do not have, most notably entanglement channel-based operations and the fact that measurement is not destructive. *PBP also can be used for traditional SIMD computation.*

## Sample pbit Layer Algorithm

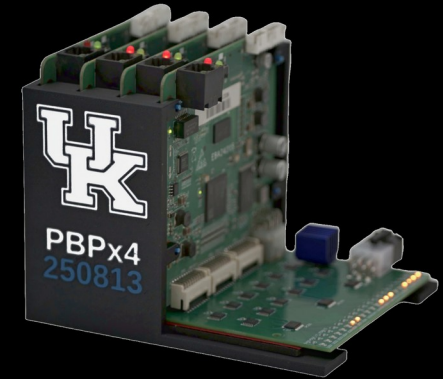
There is little point in directly using the `pbit` layer for PBP programs. However, quantum computer algorithms at the **Qubit** level can be programmed using the `pbit` layer. The following is a 4-bit ripple carry adder, adding 1 to all 4-bit values, as per Cuccaro et al, arXiv:quant-ph/0410184v1

```
void pbitripple() {
    pbit a0(0), a1(0), a2(0), a3(0);
    pbit b0(1), b1(0), b2(0), b3(0);
    pbit z(0), x(0);
    H(a0, 0); // unlike Qubits,
    H(a1, 1); // must specify groups of
    H(a2, 2); // entanglement channels
    H(a3, 3); // for Hadamard gates
    CNOT(a1,b1); CNOT(a2,b2);
    CNOT(a3,b3); CNOT(a1,x);
    CCNOT(a0,b0,x); CNOT(a2,a1);
    CCNOT(x,b1,a1); CNOT(a3,a2);
    CCNOT(a1,b2,a2); CNOT(a3,z);
    CCNOT(a2,b3,z); NOT(b1);
    NOT(b2); CNOT(x,b1);
    CNOT(a1,b2); CNOT(a2,b3);
    CCNOT(a1,b2,a2);
    CCNOT(x,b1,a1);
    CNOT(a3,a2); NOT(b2);
    CCNOT(a0,b0,x); CNOT(a2,a1);
    NOT(b1); CNOT(a1,x);
    CNOT(a0,b0); CNOT(a1,b1);
    CNOT(a2,b2); CNOT(a3,b3);
    SETMEAS(); // pick random channel
    printf("a=%d b=%d\n",
        MEAS(a0)+(MEAS(a1)<<1) +
        (MEAS(a2)<<2)+(MEAS(a3)<<3) ,
        MEAS(b0)+(MEAS(b1)<<1) +
        (MEAS(b2)<<2)+(MEAS(b3)<<3) );
}
```

## PBP References (oldest & FPGA)

H. Dietz, "How Low Can You Go?," In: Rauchwerger, L. (eds) Languages and Compilers for Parallel Computing. LCPC 2017. Lecture Notes in Computer Science(), vol 11403. Springer. 10.1007/978-3-030-35225-7\_8

H. Dietz, P. Eberhart and A. Rule, "Basic Operations And Structure Of An FPGA Accelerator For Parallel Bit Pattern Computation," 2021 International Conference on Rebooting Computing (ICRC), 2021, pp. 129-133. 10.1109/ICRC53822.2021.00029



## Parallel Bit Pattern Computing

C++ Library version 251114

<http://aggregate.org/PBP>

Professor Henry (Hank) Dietz  
Electrical and Computer Engineering Department  
University of Kentucky  
Lexington, KY 40506-0046  
[hankd@engr.uky.edu](mailto:hankd@engr.uky.edu)

Parallel bit pattern computing is a quantum-inspired model of computation. **Superposition** and ***n*-way entanglement** are modeled by each `pbit` (pattern bit) having an ordered set of  $2^n$  single-bit values. Each position in the ordered set is an **entanglement channel**. E.g., the 2-way entangled `pbit` values  $\{0,1,1,1\}$  and  $\{0,1,0,1\}$  could represent  $\{0,3,2,3\}$ , with probabilities of 25% 0, 25% 1, and 50% 3. These ordered bit sets are not directly stored, but encode as compressed patterns, with duplicate sub-patterns factored. Applicative caching avoids recomputation of sub-pattern operations. Overall, PBP can exponentially reduce both memory footprint and total number of gate-level operations.

Unlike quantum systems, users are encouraged to program parallel bit pattern computations at a relatively high level. This **CC BY 4.0** C++ library provides automatically-managed pattern bits (`pbit`) and variable-precision integer (`pint`) layers. Compiler optimizations are applied dynamically at runtime to further simplify the bit-level operations.

## pint Layer

A pattern integer, or **pint**, is an array of 1-32 **pbit** treated as a signed/unsigned integer. The precision and signedness of **pint** are variable at runtime, so that the minimum possible number of bits are active.

The usual C/C++ operators work as expected on **pint** values. Simple assignments between **int** and **pint** convert; other conversions must be explicit.

- **pint()**, **pint(v)**, **pint(v,p)**  
Create a **pint** initialized to an integer value: NaN, the **int** value **v**, or **v** with precision **p**
- **H(w)**, **H(w,m)**  
Create a **pint** **Hadamard** pattern **w** ways entangled using entanglement channels specified by mask **m**
- **p.Valid()**  
True *iff* **pint p** has a valid value (is not NaN)
- **p.Minimize()**  
Create value of **p** with fewest **pbit** possible
- **p.Extend(b)**  
Create value of **p** with **b** **pbit** precision
- **p.Promote(q)**  
Create value of **p** with minimum **pbit** precision that covers both **p** and **q** values and signedness
- **p.Logic()**  
Create **pint** with single **pbit** logic value of **p**
- **p.Rot(e)**  
Create value of **p** rotated by **e** entanglement channels (a simple phase shift)
- **p.Reset(e)**, **p.Set(e)**  
Create value of **p** with entanglement channel **e** reset or set
- **p.Dom(e)**  
Create value of **p** with bits dominoed (inverted) from entanglement channel **e** downward
- **p.Meas(e)**, **p.Meas()**, **i=p**  
Create **int** value of **p** from entanglement channel **e** or a random sample
- **p.First()**  
Create **int** value of first entanglement channel in **p** that holds a 1; returns  $2^{\text{ways}}$  if none
- **p.Ones()**  
Create **int** value number of entanglement channels in **p** that holds a 1

- **p.Min(q)**, **p.Max(q)**  
Create **pint** with minimum/maximum value from **p** or **q** for each entanglement channel
- **p.Abs()**  
Create **pint** with absolute value of **p**
- **p.Signed()**, **p.UnSigned()**  
Create **pint** forcing signed/unsigned interpretation of the **pbits** in **p**
- **p.Mul(q)**, **p.Mul(q,b)**  
Create **pint** product of **p** and **q**, but limit result precision to **b** **pbits** to save effort
- **p.Any()**, **p.All()**  
Create **int** value that is 1 *iff* any/all entanglement channels in **p** are non-zero
- **p.Summary()**, **p.Show()**  
Print debugging info for **pint p** value: either **pbit** summary or complete bit patterns

## pbit Layer

A pattern bit, or **pbit**, is logically a vector of  $2^{\text{ways}}$  bits, but is generally stored and operated upon in a heavily compressed form – a 32-way entangled **pbit** can take as little as 16 bits of storage space. A **pbit** is similar to a **Qubit** in a quantum computer, but **pbit** values are automatically allocated, maintain their value forever, and allow arbitrary fan-out; thus, they are not restricted to reversible gate operations. The basic operations include:

- **pbit()**, **pbit(v)**  
Create a **pbit** initialized to NaN or **pbit** register **v**: 0 is 0, 1 is 1, 2 is **H0**, 3 is **H1**, etc.
- **p.Valid()**  
True *iff* **pbit p** has a valid value (is not NaN)
- **p.And(q)**, **p.Or(q)**, **p.Xor(q)**, **p.Not()**  
Bitwise AND, OR, XOR, and NOT
- **p.Rot(e)**, **p.Flip(a)**  
Phase operators; **p** rotated by **e** entanglement channels or dimensionally flipped on **a**
- **p.Reset(e)**, **p.Set(e)**, **p.Tog(e)**  
Create value of **p** with entanglement channel **e** reset, set, or toggled
- **p.Dom(e)**  
Create value of **p** with bits dominoed (inverted) from entanglement channel **e** downward

- **p.Meas(e)**, **p.Meas()**  
Create **int** 0/1 value of **p** from entanglement channel **e** or a random sample
- **p.First()**  
Create **int** value of first entanglement channel in **p** that holds a 1; returns  $2^{\text{ways}}$  if none
- **p.Ones()**  
Create **int** value number of entanglement channels in **p** that holds a 1
- **p.Any()**, **p.All()**  
Create **pbit** value that is 1 *iff* any/all entanglement channels in **p** are non-zero
- **p.Show()**  
Print debugging info for **pbit p** value: complete bit patterns

The following **pbit** operations are provided solely for porting Qubit-level quantum algorithms:

- **NOT(q)**  
**Pauli X** gate; replaces **q** with  $\sim q$
- **CNOT(c,t)**  
Controlled not gate; where **c**, replaces **t** with  $\sim t$
- **CCNOT(a,b,c)**  
**Toffoli** gate; where **a** and **b**, replaces **c** with  $\sim c$
- **SWAP(i0,i1)**  
Swap values of **i0** and **i1**
- **CSWAP(c,i0,i1)**  
**Fredkin** gate; where **c**, swap **i0** and **i1**
- **H(q,c)**  
**Hadamard** gate; replaces **q** with  $q \wedge$  **Hadamard** entanglement pattern **c**
- **SETMEAS()** and **SETMEAS(m)**  
Set measurement of **rand()** channel or **m**
- **MEAS(q)**  
Measure and collapse state of **q**, returns 0/1

## RE, AC, and AoB Layers

The Regular Expression, Applicative Caching, and Array-of-Bits layers are not described here; they are considered internal, and may be changed without notice. Although the **pint** and **pbit** layers dramatically reduce gate-level operations per computation, these lower layers provide up to exponential reduction in both gate operations and in storage requirements. Performance of these layers can be summarized by calling **re.Stats()**.

# The `pbit` Layer

- `pbit p` ; allocates `p` and initializes to NaN
- `pbit p(0)` ; allocates `p` and initializes to 0
- `NOT(p)` ; applies Pauli X to `p`
- `CNOT(c, t)` ; where `c, t = ~t`
- `CCNOT(a, b, c)` ; where `a AND b, c = ~c`
- `SWAP(i0, i1)` ; swap values of `i0` and `i1`
- `CSWAP(c, i0, i1)` ; where `c`, swap values
- `H(p, c)` ; replace `p` with `p^H` pattern `c`
- `SETMEAS(m)` ; measure entanglement chan. `m`
- `MEAS(p)` ; measure chan. `m` of `p`, collapse `p`

# The `pbit` Layer

- `p.Valid()` checks for not NaN
- `p.And(q)`, `p.Or(q)`, `p.Xor(q)`, `p.Not()`
- `p.Rot(e)` rotates `e` channels
- `p.Flip(a)` flips dimensions spec'd by `a`
- `p.Reset(e)`, `p.Set(e)`, `p.Tog(e)`
- `p.Dom(e)` domino channels from `e` downward
- `p.Meas(e)`, `p.Meas()`
- `p.First()` lowest channel with value 1
- `p.Ones()` population count of 1 values in `p`
- `p.Any()`, `p.All()`, `p.Show()`

# 4-bit Quantum Adder

```
void
pbitripple()
{
  // 4-bit wide pbitripple-carry adder per Cuccaro et al arXiv:quant-ph/0410184v1
  pbit a0(0), a1(0), a2(0), a3(0), b0(1), b1(0), b2(0), b3(0), z(0), x(0);
  H(a0, 0); H(a1, 1); H(a2, 2); H(a3, 3);
  CNOT(a1,b1); CNOT(a2,b2);
  CNOT(a3,b3); CNOT(a1,x);
  CCNOT(a0,b0,x); CNOT(a2,a1);
  CCNOT(x,b1,a1); CNOT(a3,a2);
  CCNOT(a1,b2,a2); CNOT(a3,z);
  CCNOT(a2,b3,z); NOT(b1);
  NOT(b2); CNOT(x,b1);
  CNOT(a1,b2); CNOT(a2,b3);
  CCNOT(a1,b2,a2);
  CCNOT(x,b1,a1);
  CNOT(a3,a2); NOT(b2);
  CCNOT(a0,b0,x); CNOT(a2,a1);
  NOT(b1); CNOT(a1,x);
  CNOT(a0,b0); CNOT(a1,b1);
  CNOT(a2,b2); CNOT(a3,b3);
  SETMEAS();
  printf("a=%d b=%d\n",
         MEAS(a0) + (MEAS(a1) << 1) + (MEAS(a2) << 2) + (MEAS(a3) << 3),
         MEAS(b0) + (MEAS(b1) << 1) + (MEAS(b2) << 2) + (MEAS(b3) << 3));
}
```

# Notes

- **Phase kickback** doesn't happen here for the same reason that measurement need not collapse a superposition... but PBP does have phase operations
- In a quantum computer, measurement is of a random entanglement channel; here, the channel is selected as **m** in **SETMEAS (m)** ; and should be consistently used after selection

# Unphased

- How do we implement Deutsch's Algorithm without **Phase kickback**?
- E.g., use **Ones ()** to count 1s in superposition; same count as  $H(0)$  means balanced

```
pbit x(0); H(0, x);  
pbit y(0); H(1, y);  
CNOT(x, y);  
cout << (y.Ones() == x.Ones()) << "\n";
```

# The `pint` Layer

- A `pint` is automatically dynamically sized
  - Superposition with  $1 \ll \text{REWAYS}$  int values
  - Precision is  $1.. \text{REWAYS}$  pbits
  - Signed, but only magnitude if non-negative
- `p.Minimize()` minimizes pbits for `p`'s value
- `p.Logic()` makes `p` have just 1 pbit
- `p.Extend(b)` makes `p` have `b` pbits
- `p.Promote(q)` makes `p` big enough to hold `q`
- `p.Valid()` false if NaN

# The `pint` Layer

- All the usual C++ `int` operators are overloaded
- `H(w)`, `H(w, mask)` create `w`-way H pattern
- `p.Rot(e)`, `p.Flip(a)`, etc. as for `pbit`
- `p.Min(q)`, `p.Max(q)` min/max in each channel
- `p.Abs(q)` make unsigned value
- `p.Signed(e)`, `p.UnSigned()` treat pbits as...
- `p.Mul(q)` the multiply used for `*`
- `p.Mul(q, b)` multiply constrained to `b` pbits out
- `p.Summary()`, `p.Show()`

# The `pint` Layer

- `p.ReduceOp()` reduces superposition to scalar
- `p.ScanOp()` parallel-prefix of superposition
- `p.Scatter(ip, n, e)` converts an array into entangled superposition
- `p.Cover(min, max, dim)` make superposition that covers the specified range
- `p.Range(min, max, dim)` make superposition that tightly covers, with 0 padding
- `p.Gather(ip, n)` converts into an array
- `p.SortUp()`, `p.SortDown()` bitonic sort

# `pint` Square Root

- Square root by squaring...  
For example, `pintsqrt(169)` is 13

```
void
pintsqrt(int val)
{
    // Compute square root of val
    pint a(val); // 8-bit number
    pint b = pint(0).Had(4); // possible sqrts
    pint c = (b * b); // square them
    pint d = (c == a); // which were 169?
    int pos = d.First(); // first non-0 is answer
    printf("Square root of %d is %d\n", val, pos);
}
```

# pint Prime Factor

- Prime factoring by multiply...

For example, `pintfactor(143)` is 11 and 13

```
void
pintfactor(int val)
{
    // Factor val
    pint a(val); // 8-bit number
    pint b = pint(0).Had(4); // 4-bit possible 1st factor
    pint c = pint(0).Had(4,4); // 4-bit possible 2nd factor
    pint d = b * c; // multiply 'em
    pint e = (d == a); // which were 143?
    pint f = e * b;
    int spot = f.First(); // factors
    int one = c.Meas(spot);
    int two = b.Meas(spot);
    printf("%d, %d are factors of %d\n", one, two, val);
}
```

# The `pfloat` Layer

- **Buggy!**
- A `pfloat` contains three `pint` values:
  - Sign is always a single `pbit`
  - Exp is allowed to grow/shrink as needed
  - Frac with choice of two normalization rules:  
to MSB or to LSB
- Standard C++ `float` operators and library functions (e.g., `p.ArcTan()`)
- All the `pint` extensions

# $\pi$ in C using `int` sampling

```
#include <stdlib.h>
#include <stdio.h>

int main(int argc, char **argv) {
    int i, j, intervals = atoi(argv[1]);
    int w = intervals * intervals, sum = 0;
    for (i=0; i<intervals; ++i) {
        int h = w / (intervals + ((i*i) / intervals));
        for (j=0; j<intervals; ++j) {
            if (h > j) ++sum;
        }
    }
    printf("Pi is roughly %f\n", (4.0 * sum) / w);
    return(0);
}
```

# $\pi$ in PBP, pint sampling

```
#include "pbp.h" // PBP classes and support

int main(int argc, char **argv) {
    int bits = atoi(argv[1]); // number of pbits
    pint intervals(1 << bits); // intervals in pbits
    pint w(1 << (2 * bits)); // int scaling factor
    pint x = pint(0).Had(bits); // all x values
    pint y = pint(0).Had(bits,bits); // all y values
    pint h = w / (((x * x) >> bits) + intervals);
    pint r = (h > y); // r is 1 where below curve
    // count 1s; quantum would sample probability
    double pi = (4.0 * r.Pop()) / (1 << REWAYS);
    printf("Pi is roughly %f\n", pi);
    return(0);
}
```